

# Tropical Optimization Framework for Analytical Hierarchy Process

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LMS Workshop on  
Tropical Mathematics & its Applications  
University of Birmingham, November 15, 2017

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# Introduction: Tropical Optimization

- ▶ **Tropical (idempotent) mathematics** focuses on the theory and applications of semirings with idempotent addition
- ▶ The **tropical optimization problems** are formulated and solved within the framework of tropical mathematics
- ▶ Many problems have **objective functions** defined on vectors over **idempotent semifields** (semirings with multiplicative inverses)
- ▶ The problems find **applications** in many areas to provide new efficient solutions to various old and novel problems in
  - ▶ project scheduling,
  - ▶ location analysis,
  - ▶ transportation networks,
  - ▶ decision making,
  - ▶ discrete event systems

# Idempotent Algebra: Definitions and Notation

## Idempotent Semifield

- ▶ *Idempotent semifield*: the algebraic system  $\langle \mathbb{X}, 0, \mathbb{1}, \oplus, \otimes \rangle$
- ▶ The binary operations  $\oplus$  and  $\otimes$  are *associative and commutative*
- ▶ The carrier set  $\mathbb{X}$  has neutral elements, *zero*  $0$  and *identity*  $\mathbb{1}$
- ▶ Multiplication  $\otimes$  is *distributive* over addition
- ▶ Addition  $\oplus$  is *idempotent*:  $x \oplus x = x$  for all  $x \in \mathbb{X}$
- ▶ Multiplication  $\otimes$  is *invertible*: for each nonzero  $x \in \mathbb{X}$ , there exists an inverse  $x^{-1} \in \mathbb{X}$  such that  $x \otimes x^{-1} = \mathbb{1}$
- ▶ *Algebraic completeness*: the equation  $x^p = a$  is solvable for any  $a \in \mathbb{X}$  and integer  $p$  (there exist powers with rational exponents)
- ▶ *Notational convention*: the multiplication sign  $\otimes$  will be omitted

## Semifield $\mathbb{R}_{\max, \times}$ (Max-Algebra)

- ▶ **Definition:**  $\mathbb{R}_{\max, \times} = \langle \mathbb{R}_+, 0, 1, \max, \times \rangle$  with  $\mathbb{R}_+ = \{x \in \mathbb{R} | x \geq 0\}$
- ▶ **Carrier set:**  $\mathbb{X} = \mathbb{R}_+$ ; **zero and identity:**  $0 = 0, 1 = 1$
- ▶ **Binary operations:**  $\oplus = \max$  and  $\otimes = \times$
- ▶ **Idempotent addition:**  $x \oplus x = \max(x, x) = x$  for all  $x \in \mathbb{R}_+$
- ▶ **Multiplicative inverse:** for each  $x \in \mathbb{R}_+ \setminus \{0\}$ , there exists  $x^{-1}$
- ▶ **Power notation:**  $x^y$  is routinely defined for each  $x, y \in \mathbb{R}_+$
- ▶ Further examples of real idempotent semifields:

$$\mathbb{R}_{\max, +} = \langle \mathbb{R} \cup \{-\infty\}, -\infty, 0, \max, + \rangle,$$

$$\mathbb{R}_{\min, +} = \langle \mathbb{R} \cup \{+\infty\}, +\infty, 0, \min, + \rangle,$$

$$\mathbb{R}_{\min, \times} = \langle \mathbb{R}_+ \cup \{+\infty\}, +\infty, 1, \min, \times \rangle$$

## Vector and Matrix Algebra Over $\mathbb{R}_{\max, \times}$

- ▶ The scalar idempotent semifield  $\mathbb{R}_{\max, \times}$  is routinely extended to idempotent systems of vectors in  $\mathbb{R}_+^n$  and of matrices in  $\mathbb{R}_+^{m \times n}$
- ▶ The matrix and vector operations follow the standard entry-wise formulas with the addition  $\oplus = \max$  and the multiplication  $\otimes = \times$
- ▶ For any vectors  $\mathbf{a} = (a_i)$  and  $\mathbf{b} = (b_i)$  in  $\mathbb{R}_+^n$ , and a scalar  $x \in \mathbb{R}_+$ , the vector operations follow the conventional rules

$$\{\mathbf{a} \oplus \mathbf{b}\}_i = a_i \oplus b_i, \quad \{x\mathbf{a}\}_i = xa_i$$

- ▶ For any matrices  $\mathbf{A} = (a_{ij}) \in \mathbb{R}_+^{m \times n}$ ,  $\mathbf{B} = (b_{ij}) \in \mathbb{R}_+^{m \times n}$  and  $\mathbf{C} = (c_{ij}) \in \mathbb{R}_+^{n \times l}$ , and  $x \in \mathbb{R}_+$ , the matrix operations are given by

$$\{\mathbf{A} \oplus \mathbf{B}\}_{ij} = a_{ij} \oplus b_{ij}, \quad \{\mathbf{AC}\}_{ij} = \bigoplus_{k=1}^n a_{ik}c_{kj}, \quad \{x\mathbf{A}\}_{ij} = xa_{ij}$$

## Vector and Matrix Algebra Over $\mathbb{R}_{\max, \times}$

- ▶ All vectors are column vectors, unless otherwise specified
- ▶ The *zero vector* and *vector of ones*:  
 $\mathbf{0} = (0, \dots, 0)^T$ ,  $\mathbf{1} = (1, \dots, 1)^T$
- ▶ *Multiplicative conjugate transposition* of a nonzero column vector  $\mathbf{x} = (x_i)$  is the row vector  $\mathbf{x}^- = (x_i^-)$ , where  $x_i^- = x_i^{-1}$  if  $x_i \neq 0$ , and  $x_i^- = 0$  otherwise
- ▶ The *zero matrix* and *identity matrix*:  
 $\mathbf{0} = (0)$ ,  $\mathbf{I} = \text{diag}(1, \dots, 1)$
- ▶ *Multiplicative conjugate transposition* of a nonzero matrix  $\mathbf{A} = (a_{ij})$  is the matrix  $\mathbf{A}^- = (a_{ij}^-)$ , where  $a_{ij}^- = a_{ji}^{-1}$  if  $a_{ji} \neq 0$ , and  $a_{ij}^- = 0$  otherwise
- ▶ *Integer powers* of square matrices:

$$\mathbf{A}^0 = \mathbf{I}, \quad \mathbf{A}^p = \mathbf{A}^{p-1} \mathbf{A} = \mathbf{A} \mathbf{A}^{p-1}, \quad p \geq 1$$

## Square Matrices

- ▶ *Trace*: the trace of a matrix  $\mathbf{A} = (a_{ij}) \in \mathbb{R}_+^{n \times n}$  is given by

$$\text{tr } \mathbf{A} = a_{11} \oplus \cdots \oplus a_{nn}$$

- ▶ *Eigenvalue*: a scalar  $\lambda$  such that there is a vector  $x \neq 0$  to satisfy

$$\mathbf{A}x = \lambda x$$

- ▶ *Spectral radius*: the maximum eigenvalue given by

$$\rho = \text{tr } \mathbf{A} \oplus \cdots \oplus \text{tr}^{1/m}(\mathbf{A}^m)$$

- ▶ *Asterate*: the asterate operator (the Kleene star) is given by

$$\mathbf{A}^* = \mathbf{I} \oplus \mathbf{A} \oplus \cdots \oplus \mathbf{A}^{n-1}, \quad \rho \leq 1$$



# Tropical Optimization Problems: Solution Examples

## Problem with Pseudo-Quadratic Objective

Given a matrix  $A \in \mathbb{R}_+^{n \times n}$ , find positive vectors  $x \in \mathbb{R}_+^n$  that solve the problem

$$\min_{x > 0} x^- A x$$

## Theorem

Let  $A$  be a matrix with tropical spectral radius  $\lambda > 0$ , and denote  $B = (\lambda^{-1} A)^*$ . Then,

- ▶ the minimum of  $x^- A x$  is equal to  $\lambda$ ;
- ▶ all positive solutions are given by

$$x = B u, \quad u > 0$$

## Maximization Problem with Hilbert (Range, Span) Seminorm

Given a matrix  $B \in \mathbb{R}_+^{n \times m}$ , find positive vectors  $u \in \mathbb{R}_+^l$  that solve the problems

$$\max_{u > 0} \mathbf{1}^T B u (B u)^{-1} \qquad \min_{u > 0} \mathbf{1}^T B u (B u)^{-1}$$

### Lemma

Let  $B$  be a positive matrix, and  $B_{lk}$  be the matrix derived from  $B = (b_k)_{k=1}^m$  by fixing the entry  $b_{lk}$  and replacing the others by 0.

- ▶ The **maximum** of  $\mathbf{1}^T B u (B u)^{-1}$  is equal to  $\Delta = \mathbf{1}^T B B^{-1}$
- ▶ All positive solutions are given by

$$u = (I \oplus B_{lk}^- B) v, \quad v > 0,$$

where the indices  $k$  and  $l$  satisfy the condition  $\mathbf{1}^T b_k b_{lk}^{-1} = \Delta$

## Minimization Problem with Hilbert (Range, Span) Seminorm

### Lemma

Let  $B$  be a matrix without zero rows and columns.

- ▶ The **minimum** of  $\mathbf{1}^T B u (B u)^{-1} \mathbf{1}$  is equal to  $\Delta = (B(\mathbf{1}^T B)^{-})^{-1} \mathbf{1}$ .
- ▶ Denote by  $\widehat{B}$  be the sparsified matrix with entries:

$$\widehat{b}_{ij} = \begin{cases} 0, & \text{if } b_{ij} < \Delta^{-1} \mathbf{1}^T \mathbf{b}_j; \\ b_{ij}, & \text{otherwise.} \end{cases}$$

Let  $\mathcal{B}$  be the set of matrices obtained from  $\widehat{B}$  by fixing one nonzero entry in each row and setting the others to 0.

Then, all positive solutions are given by

$$u = (I \oplus \Delta^{-1} B_1^{-1} \mathbf{1}^T B) v, \quad v > 0, \quad B_1 \in \mathcal{B}$$

# Analytical Hierarchy Process: Traditional Approach

## Pairwise Comparison

- ▶ Given  $m$  criteria and  $n$  choices, the problem is to find priorities of choices from pairwise comparisons of criteria and of choices
- ▶ Outcome of comparison is given by a matrix  $A = (a_{ij})$ , where  $a_{ij}$  shows the relative priority of choice  $i$  over  $j$
- ▶ Note that  $a_{ij} = 1/a_{ji} > 0$
- ▶ Scale (Saaty, 2005):

$a_{ij}$	Meaning
1	$i$ equally important as $j$
3	$i$ moderately more important than $j$
5	$i$ strongly more important than $j$
7	$i$ very strongly more important than $j$
9	$i$ extremely more important than $j$

## Consistency

- ▶ A pairwise comparison matrix  $A$  is consistent if its entries are transitive to satisfy the condition  $a_{ij} = a_{ik}a_{kj}$  for all  $i, j, k$
- ▶ Each consistent matrix  $A$  has unit rank and is given by  $A = xx^T$ , where  $x$  is a positive vector that entirely specifies  $A$
- ▶ If a comparison matrix  $A$  is consistent, the vector  $x$  represents, up to a positive factor, the individual priorities of choices
- ▶ Since the comparison matrices are usually inconsistent, a problem arises to approximate these matrices by consistent matrices

## Principal Eigenvector Method and Weighted Sum Solution

- ▶ The traditional AHP uses approximation of pairwise comparison matrices by consistent matrices with the principal eigenvectors
- ▶ Let  $A_0$  be a matrix of pairwise comparison of criteria, and  $A_k$  be a matrix of pairwise comparison of choices for criterion  $k$
- ▶ Let  $w = (w_k)_{k=1}^m$  be the principal eigenvector of  $A_0$ : the vector of priorities (weights) for criteria
- ▶ Let  $x_k$  be the principal eigenvector of  $A_k$ : the vector of priorities of choices with respect to criterion  $k$
- ▶ The resulting vector  $x$  of priorities of choices is calculated as

$$\mathbf{x} = \sum_{k=1}^m w_k \mathbf{x}_k$$

# Minimax approximation based AHP

## Log-Chebyshev Approximation of Comparison Matrices

- Consider the problem to approximate a pairwise comparison matrix  $A = (a_{ij})$  by a consistent matrix  $X = (x_{ij})$ , where

$$a_{ij} = 1/a_{ji}, \quad x_{ij} = x_i/x_j$$

- The log-Chebyshev distance between  $A$  and  $X$  is defined as

$$\max_{1 \leq i, j \leq n} |\log a_{ij} - \log x_{ij}| = \log \max_{1 \leq i, j \leq n} \max \left( \frac{a_{ij}}{x_{ij}}, \frac{x_{ij}}{a_{ij}} \right)$$

- Minimizing the log-Chebyshev distance is equivalent to minimizing

$$\max_{1 \leq i, j \leq n} \max \left( \frac{a_{ij}}{x_{ij}}, \frac{x_{ij}}{a_{ij}} \right) = \max_{1 \leq i, j \leq n} \max \left( \frac{a_{ij}x_j}{x_i}, \frac{a_{ji}x_i}{x_j} \right) = \max_{1 \leq i, j \leq n} \frac{a_{ij}x_j}{x_i}$$

## Approximation as Tropical Optimization Problem

- ▶ Tropical representation of the objective function in terms of  $\mathbb{R}_{\max, \times}$

$$\max_{1 \leq i, j \leq n} \frac{a_{ij} x_j}{x_i} = \bigoplus_{1 \leq i, j \leq n} x_i^{-1} a_{ij} x_j = \mathbf{x}^{-} \mathbf{A} \mathbf{x}$$

- ▶ In the framework of the idempotent semifield  $\mathbb{R}_{\max, \times}$ , the minimax approximation problem takes the form

$$\min_{\mathbf{x} > \mathbf{0}} \mathbf{x}^{-} \mathbf{A} \mathbf{x}$$

### Theorem

Let  $\mathbf{A}$  be a pairwise comparison matrix with tropical spectral radius  $\lambda$ , and  $\mathbf{B} = (\lambda^{-1} \mathbf{A})^*$ . Then,

- ▶ all priority vectors are given by

$$\mathbf{x} = \mathbf{B} \mathbf{u}, \quad \mathbf{u} > \mathbf{0}$$



## Weighted Approximation Under Several Criteria

- ▶ Simultaneous minimax approximation of the matrices  $A_k = (a_{ij}^{(k)})$  with weights  $w_k > 0$  by a consistent matrix involves minimizing

$$\max_{1 \leq k \leq m} w_k \left( \max_{1 \leq i, j \leq n} (a_{ij}^{(k)} x_j) / x_i \right) = \max_{1 \leq i, j \leq n} \max_{1 \leq k \leq m} (w_k a_{ij}^{(k)}) x_j / x_i.$$

- ▶ In terms of  $\mathbb{R}_{\max, \times}$ , the approximation problem takes the form

$$\min_{x > 0} x^- (w_1 A_1 \oplus \cdots \oplus w_m A_m) x$$

### Theorem

Let  $A_1, \dots, A_m$  be comparison matrices,  $w_1, \dots, w_m$  be weights,  $C = w_1 A_1 \oplus \cdots \oplus w_m A_m$  be a matrix with tropical spectral radius  $\mu$ , and  $B = (\mu^{-1} C)^*$ . Then,

- ▶ all priority vectors are given by

$$x = Bu, \quad u > 0$$

## Most and Least Differentiating Priority Vectors

- ▶ The priority vectors  $x = Bu$  obtained by the minimax approximation may be not unique up to a positive factor
- ▶ Further analysis is then needed to reduce to a few representative solutions, such as some “best” and “worst” priority vectors
- ▶ One can take two vectors that most and least differentiate between the choices with the highest and lowest priorities
- ▶ The most and least differentiating priority vectors are obtained by maximizing and minimizing the contrast ratio

$$\max_{1 \leq i \leq n} x_i / \min_{1 \leq i \leq n} x_i = \max_{1 \leq i \leq n} x_i \times \max_{1 \leq i \leq n} x_i^{-1}$$

- ▶ In terms of the semifield  $\mathbb{R}_{\max, \times}$ , the contrast ratio is written as

$$\mathbf{1}^T \mathbf{x} \mathbf{x}^{-1} = \mathbf{1}^T \mathbf{B} \mathbf{u} (\mathbf{B} \mathbf{u})^{-1}$$

## Most Differentiating Priority Vector

- ▶ In terms of the semifield  $\mathbb{R}_{\max, \times}$ , the problem to maximize the contrast ratio is written as

$$\max_{\mathbf{u} > \mathbf{0}} \mathbf{1}^T \mathbf{B} \mathbf{u} (\mathbf{B} \mathbf{u})^{-1} \mathbf{1}$$

- ▶ If  $\mathbf{u}$  is a solution, then  $\mathbf{x} = \mathbf{B} \mathbf{u}$  is the most differentiating vector

## Most Differentiating Priority Vector

### Theorem

Let  $B$  be a matrix defining a set of priority vectors  $x = Bu$ ,  $u > 0$ , and  $B_{lk}$  be the matrix obtained from  $B = (b_j)$  by fixing the entry  $b_{lk}$  and replacing the others by 0.

- ▶ The **maximum** of  $\mathbf{1}^T Bu (Bu)^{-1}$  is equal to  $\Delta = \mathbf{1}^T B B^{-1}$
- ▶ The **most differentiating priority vectors** are given by

$$x = B(I \oplus B_{lk}^- B)v, \quad v > 0,$$

where the indices  $k$  and  $l$  satisfy the condition  $\mathbf{1}^T b_k b_{lk}^{-1} = \Delta$

## Least Differentiating Priority Vector

- ▶ In terms of the semifield  $\mathbb{R}_{\max, \times}$ , the problem to minimize the contrast ratio is written as

$$\min_{\mathbf{u} > \mathbf{0}} \mathbf{1}^T \mathbf{B} \mathbf{u} (\mathbf{B} \mathbf{u})^{-1} \mathbf{1}$$

- ▶ If  $\mathbf{u}$  is a solution, then  $\mathbf{x} = \mathbf{B} \mathbf{u}$  is the least differentiating vector

## Least Differentiating Priority Vector

### Theorem

Let  $B$  be a matrix defining a set of priority vectors  $x = Bu$ ,  $u > 0$ .

- ▶ The **minimum** of  $\mathbf{1}^T B u (B u)^{-1} \mathbf{1}$  is equal to  $\Delta = (B(\mathbf{1}^T B)^{-1})^{-1}$ .
- ▶ Denote by  $\hat{B}$  be the sparsified matrix with entries:

$$\hat{b}_{ij} = \begin{cases} 0, & \text{if } b_{ij} < \Delta^{-1} \mathbf{1}^T \mathbf{b}_j; \\ b_{ij}, & \text{otherwise.} \end{cases}$$

Let  $\mathcal{B}$  be the set of matrices obtained from  $\hat{B}$  by fixing one nonzero entry in each row and setting the others to 0.

Then the **least differentiating priority vectors** are given by

$$u = (I \oplus \Delta^{-1} B_1^{-1} \mathbf{1} \mathbf{1}^T B) v, \quad v > 0, \quad B_1 \in \mathcal{B}$$

# Illustrative Example: Selecting Plan for Vacation

## Problem: Select a Place to Spend a Week (Saaty, 1977)

- Five criteria: (1) cost of the trip from Philadelphia, (2) sight-seeing opportunities, (3) entertainment (doing things), (4) way of travel, (5) eating places; with the criteria comparison matrix

$$A_0 = \begin{pmatrix} 1 & 1/5 & 1/5 & 1 & 1/3 \\ 5 & 1 & 1/5 & 1/5 & 1 \\ 5 & 5 & 1 & 1/5 & 1 \\ 1 & 5 & 5 & 1 & 5 \\ 3 & 1 & 1 & 1/5 & 1 \end{pmatrix}$$

- Four places: (1) short trips from Philadelphia (i.e., New York, Washington, Atlantic City, New Hope, etc.), (2) Quebec, (3) Denver, (4) California

## Problem (cont.)

- Pairwise comparison matrices of places with respect to criteria

$$\mathbf{A}_1 = \begin{pmatrix} 1 & 3 & 7 & 9 \\ 1/3 & 1 & 6 & 7 \\ 1/7 & 1/6 & 1 & 3 \\ 1/9 & 1/7 & 1/3 & 1 \end{pmatrix}, \quad \mathbf{A}_2 = \begin{pmatrix} 1 & 1/5 & 1/6 & 1/4 \\ 5 & 1 & 2 & 4 \\ 6 & 1/2 & 1 & 6 \\ 4 & 1/4 & 1/6 & 1 \end{pmatrix},$$

$$\mathbf{A}_3 = \begin{pmatrix} 1 & 7 & 7 & 1/2 \\ 1/7 & 1 & 1 & 1/7 \\ 1/7 & 1 & 1 & 1/7 \\ 2 & 7 & 7 & 1 \end{pmatrix}, \quad \mathbf{A}_4 = \begin{pmatrix} 1 & 4 & 1/4 & 1/3 \\ 1/4 & 1 & 1/2 & 3 \\ 4 & 2 & 1 & 3 \\ 3 & 1/3 & 1/3 & 1 \end{pmatrix},$$

$$\mathbf{A}_5 = \begin{pmatrix} 1 & 1 & 7 & 4 \\ 1 & 1 & 6 & 3 \\ 1/7 & 1/6 & 1 & 1/4 \\ 1/4 & 1/3 & 4 & 1 \end{pmatrix}$$



## Solution: Evaluating Priority Vector (Weights) for Criteria

- ▶ The tropical spectral radius of the comparison matrix  $\mathbf{A}_0$

$$\lambda = (a_{14}a_{43}a_{32}a_{21})^{1/4} = 5^{3/4}$$

- ▶ The Kleene star of the matrix  $\lambda^{-1}\mathbf{A}_0$

$$\begin{aligned}
 (\lambda^{-1}\mathbf{A}_0)^* &= \mathbf{I} \oplus \lambda^{-1}\mathbf{A}_0 \oplus \lambda^{-2}\mathbf{A}_0^2 \oplus \lambda^{-3}\mathbf{A}_0^3 \oplus \lambda^{-4}\mathbf{A}_0^4 \\
 &= \begin{pmatrix} 1 & 5^{-1/4} & 5^{-1/2} & 5^{-3/4} & 5^{-1/2} \\ 5^{1/4} & 1 & 5^{-1/4} & 5^{-1/2} & 5^{-1/4} \\ 5^{1/2} & 5^{1/4} & 1 & 5^{-1/4} & 1 \\ 5^{3/4} & 5^{1/2} & 5^{1/4} & 1 & 5^{1/4} \\ 3 \cdot 5^{-3/4} & 3 \cdot 5^{-1} & 3 \cdot 5^{-5/4} & 3 \cdot 5^{-3/2} & 3 \cdot 5^{-5/4} \end{pmatrix}
 \end{aligned}$$

- ▶ The priority (weight) vector for criteria (pseudo-quadratic problem)

$$\mathbf{w} = (1, 5^{1/4}, 5^{1/2}, 5^{3/4}, 3 \cdot 5^{-3/4})$$

## Derivation of All Priority Vectors for Places

- ▶ The weighted combination of comparison matrices of places

$$\begin{aligned}
 C &= A_1 \oplus 5^{1/4} A_2 \oplus 5^{1/2} A_3 \oplus 5^{3/4} A_4 \oplus (3 \cdot 5^{-3/4}) A_5 \\
 &= \begin{pmatrix} 5^{3/4} & 7 \cdot 5^{1/2} & 7 \cdot 5^{1/2} & 9 \\ 5^{5/4} & 5^{3/4} & 6 & 3 \cdot 5^{3/4} \\ 4 \cdot 5^{3/4} & 2 \cdot 5^{3/4} & 5^{3/4} & 3 \cdot 5^{3/4} \\ 3 \cdot 5^{3/4} & 7 \cdot 5^{1/2} & 7 \cdot 5^{1/2} & 5^{3/4} \end{pmatrix}
 \end{aligned}$$

- ▶ The tropical spectral radius of matrix  $C$

$$\mu = (c_{13}c_{31})^{1/2} = 2 \cdot 5^{5/8} 7^{1/2}$$

## Derivation of All Priority Vectors for Places (cont.)

- ▶ The Kleene star of matrix  $\mu^{-1}C$

$$\begin{aligned}
 (\mu^{-1}C)^* &= I \oplus \mu^{-1}C \oplus \mu^{-2}C^2 \oplus \mu^{-3}C^3 \\
 &= \begin{pmatrix} 1 & r/4 & r/4 & 3/4 \\ 3/r & 1 & 3/4 & 3/r \\ 4/r & 1 & 1 & 3/r \\ 1 & r/4 & r/4 & 1 \end{pmatrix}, \quad r = 2 \cdot 7^{1/2} 5^{-1/8} \approx 4.33
 \end{aligned}$$

- ▶ All solution vectors (pseudo-quadratic problem)

$$x = Bu, \quad B = \begin{pmatrix} 1 & r/4 & 3/4 \\ 3/r & 1 & 3/r \\ 4/r & 1 & 3/r \\ 1 & r/4 & 1 \end{pmatrix}, \quad u > 0$$

## Evaluation of Most Differentiating Solutions

- ▶ The vectors with maximum differentiation between choices of lowest and highest priorities (maximization of Hilbert seminorm)

$$\mathbf{x}_1 = \begin{pmatrix} 1 \\ 3/r \\ 4/r \\ 1 \end{pmatrix} u, \quad \mathbf{x}_2 = \begin{pmatrix} 3/4 \\ 3/r \\ 3/r \\ 1 \end{pmatrix} v, \quad u, v > 0, \quad r = 2 \cdot 7^{1/2} 5^{-1/8} \approx 4.33$$

- ▶ Examples of vectors with  $u = v = 1$

$$\mathbf{x}_1 \approx (1.00, 0.69, 0.92, 1.00)^T, \quad \mathbf{x}_2 \approx (0.75, 0.69, 0.69, 1.00)^T$$

- ▶ The priority order of places according to the vectors

$$(4) \equiv (1) \succ (3) \succ (2), \quad (4) \succ (1) \succ (3) \equiv (2)$$

## Evaluation of Least Differentiating Solutions

- ▶ The vectors with minimum differentiation between choices of lowest and highest priorities (minimization of Hilbert seminorm)

$$\mathbf{x}_1 = \begin{pmatrix} 1 \\ 4/r \\ 4/r \\ 1 \end{pmatrix} u, \quad u > 0, \quad r = 2 \cdot 7^{1/2} 5^{-1/8} \approx 4.33$$

- ▶ Example of vectors with  $u = 1$  and related priority order

$$\mathbf{x}_1 \approx (1.00, 0.92, 0.92, 1.00)^T, \quad (4) \equiv (1) \succ (3) \equiv (2)$$

- ▶ Combined new orders versus order by (Saaty, 1977)

$$\text{NEW: } (4) \succeq (1) \succ (3) \succeq (2) \quad \text{OLD: } (1) \succ (3) \succ (4) \succ (2)$$

# Concluding Remarks

- ▶ We have proposed a new implementation of the AHP method, based on minimax approximation and tropical optimization
- ▶ The new AHP implementation uses log-Chebyshev matrix approximation instead of the principal eigenvector method
- ▶ The weights of criteria are incorporated into the evaluation of the priorities of choices rather than used to form the result directly

# Concluding Remarks

- ▶ Since the solution obtained is usually non-unique, a technique has been proposed to find two representative priority vectors
- ▶ As such solutions, those vectors are taken which most and least differentiate between choices with the highest and lowest priorities
- ▶ The above problems have been formulated in the framework of tropical mathematics, and solved as tropical optimization problems
- ▶ Exact solutions to the problems have been given in a compact vector form ready for further analysis and practical implementation